

Using Contextual Security Policies for Threat Response

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Information systems service goals





Security is one of many adjustment variables
 Compromises are generally static at design time



But ...

→ Security is not static

- New vulnerabilities
- New users and usages
- New attackers
- Nor are the other variables
 - Reflect the evolution of the IS (new hardware & software)
 - Maintain a better balance between the different requirements
- → The compromise between these variables needs to change
 - Respond to threat
 - Dynamic security policies



Threat response system



→ Reactivity

- Automated response process
- On-time deployment of response according to threats
- On-time withdrawal of response
- → Reliability
 - Consistency of the threat characterization system (reliable alerts)
 - Relevance of selected countermeasures
 - Application of countermeasures to multiple enforcement points

Ease of use

- Ease of deployment (avoid or limit additional systems)
- Ease of countermeasures definition



How to fulfill the requirements?



- Clear identification of the threat, source and victim
- Policy-oriented approach
 - Adapt security level to the threat level (dynamic policy)
 - Compromise between security, performance, convenience, etc.
 - Avoid the deployment of additional systems
- Organization-based approach
 - Abstract vs concrete level of rules
 - Provide local reactions but responding to global constraints
- Context-based approach
 - Trigger security rules thanks to active contexts
 - In particular, threat contexts to trigger countermeasures



5

Access Control Policy (1)



Discretionary Access Control (DAC)

- Manage privileges of subjects on objects
- Definition of an access matrix to describe authorizations.

```
(Subject, Object, Privilege)
Ex: (host1, file1, rw),
Means that host1 has the privilege of read and write on file1.
```

Limitations

- Many subjects and objects to describe
- Scalability issues (definition and administration)
- Poor expressiveness (static policy)



Access Control Policy (2)



→ Role-Based Access Control (RBAC)

- Abstract subjects into roles
- Manage permissions of actions

Permission (Role, Action, Object) UA \subseteq UxR, user-to-role assignment

```
Ex: Permission(group1, read, file1), with host1 ∈ group1, Means that group1, thus host1, is permitted to read file1.
```

Limitations

- Only provides means to group subjects, but not actions and objects
- Only manages permissions (no explicit prohibition)
- Limited expressiveness of the security rules (static policy)

+ Empower (Organization, Subject, Role)

- + Consider (Organization, Action, Activity)
- + Use(Organization, Object, View)

Manage entities through organizations

Organization-Based Access Control (Or-BAC)

• Abstract *subjects* into *roles*

Access Control Policy (3)

- Abstract actions into activities
- Abstract objects into views







DIMVA06 - July 13-14, 2006

Unrestricted

Access Control Policy (3)

Organization-Based Access Control (Or-BAC)

- Manage entities through organizations
- Abstract *subjects* into *roles*
- Abstract actions into activities
- Abstract *objects* into *views*
- Provide not only *permissions*, but also *prohibitions*/obligations

Security_rule(Type, Organization, Role, Activity, View)

- + Empower(Organization, Subject, Role)
- + Consider (Organization, Action, Activity)
- + Use(Organization, Object, View)

With Type={permission, prohibition, obligation}





Unrestricted

Access Control Policy (3)

Organization-Based Access Control (Or-BAC)

- Manage entities through organizations
- Abstract subjects into roles
- Abstract actions into activities
- Abstract *objects* into *views*
- Provide not only *permissions*, but also *prohibitions*/obligations
- Trigger rules provided *contexts* (dynamic policy)

```
Security_rule(Type, Organization, Role, Activity, View, Context)
+ Empower(Organization, Subject, Role)
+ Consider(Organization, Action, Activity)
+ Use(Organization, Object, View)
+ Hold(Organization, Subject, Action, Object, Context)
```

With Type={permission, prohibition, obligation}







→ In the organization corp, the activity read mail is permitted for the role *mail user* on the view *mailserver* in a normal context.





Hold(corp, bob, tcp/110, mel1, normal)

In the organization *corp*, the context *normal* is being \rightarrow held for user *bob* making action *tcp/110* on object *mel1*.





In the organization corp, bob is a mail_user subject, tcp/110 is a read_mail action and mel1 is a mailserver object.









- assumed reasonably accurate for the purpose of this work)
- Output: IDMEF messages (Intrusion Detection Message Exchange Format)



- Input: IDMEF messages (characterized threats)
- Role: Dynamically extract new policy rules considering threats
- Output: New policy rules (or instances)



- Additional input: Policy definition
 - Generic Or-BAC policy (security rules, i.e. abstract policy)
 - Context definition (conditions to trigger contexts, i.e. *hold* predicates)
 - Context data (base of additional facts, apart from alerts, such as time, cartography, etc.)



- Input: Policy rules
- Role: Prepare the policy for local enforcement
- Output: PEP configurations



- Input: PEP configurations
- Role: Apply new configurations, i.e. enforce the policy
- Potential output: Events/alerts (PEPs acting as sensors)

From alerts to new policies (1)



→ Alerts provide identification of source, victim and threat

- IDMEF.Source: IP address, DNS name, network mask, etc.
- IDMEF.Target: IP address, DNS name, network mask, etc.
- IDMEF.Classification: Reference (ex: CVE-2005-1133)

Mapping strategy

- Trigger a hold(org, subject, action, object, context) from alerts ensuring adequate response to the threat
- Example
 - hold(corp, bob, tcp/110, mel1, pop_threat)



From alerts to new policies (2)



Derive concrete permissions/prohibitions (new policies) from security_rules and hold facts





From alerts to new policies (3)



Concrete permissions/prohibitions managed by the PDP

- Deployment: Adapt new policy instances into a concrete enforcement strategy
 - Block a port on a firewall,
 - Stop/reconfigure a service,
 - Etc.
- Translation: Adapt policy rules to PEPs type and implementation
 - Type: "A firewall rule"
 - Implementation: "A Netfilter firewall rule"
- PEPs receive new configurations by the PDP to enforce the new policy



Minimal requirements



Contexts allow expression of minimal requirements

- Ex: 3 different paths to read mail (pop, imap and webmail)
- During working hours, availability is considered more important than confidentiality and integrity for mail
- If all paths to mail are threatened, re-open webmail to fulfill availability requirement, whatever the threat

Security_rule(perm,corp,mail_user,read_webmail,mailserver,minimal)

With minimal=pop_threat&imap_threat&webmail_threat&working_hours)



Keeping a sensitive path open to maintain availability

Conclusion

is questionable?

- Availability is a crucial requirement
- Other means can be deployed to ensure confidentiality and integrity
 - In particular, responses can be defined to switch between different requirements of authentication, ciphering, etc.
 - Provisional authorizations
- → Results are encouraging
 - An implementation of the PIE/PDP in Prolog confirms the feasibility of the approach



Future work

Mapping is a great part of the work in progress

- Provide relevant hold facts to ensure adequate responses
- Scale of the response

Context lifetime

- In a first time, static context lifetime based on expertise
- Next step: characterizing the absence of threat (anti-alerts?)

→ Experiments



Questions?





